

V-T-A

a 1937 Turing Machine (one a mathematical
 def. of computation - algorithm
 idea: Turing Machine

environment  $\rightarrow$ sequence type

$A = \text{alphabet (per state)}$

Possible agent
attras

- 1) replace a symbol in n^{th} cell by another symbol
- 2) H e can change his internal state of mind (H e has a finite # of available internal states)
- 3) Make a move left, right, or stop put



$$\phi = \{a_1, \dots, a_n\} \text{ set of state}$$

A

Pla_B

Print $p_1 \otimes x_A \rightarrow A, p(a_3 \otimes 5) = 9$

State $\text{SIGKD} \rightarrow \mathcal{Q}$

мы $\Pi : \mathcal{G}K\mathcal{A} \rightarrow \Sigma^{-1}, 0, 1\}$

Recursively enumerable sets

Recursive functions

give the machine a number

Choose an alphabet $A = \{0, 1\}$

A configuration with ~~all~~ n 1's and then all 0's.

Machine starts in left most cell in state q_0

q^*
 q^*
 a^*

$$P(q^* a^*) = a^*$$

$$S(q^* a^*) = q^*$$

$$H(q^* a^*) = \emptyset$$

Suppose eventually machine stops with number m

Def: A function $f: \mathbb{N} \rightarrow \mathbb{N}$ is called partial recursive if there is a

turing machine such that when started on the tape with $\#n$ on it — — it halts with m on its tape $\leftrightarrow$
 $f(n) = m$

An number n in A will eventually find out by rat you never will.

Def: A is recursive if both A and $\mathbb{N} - A$ are rec

decidable

TM (universal machine)

subset

$$C = \{n \mid n \in W_n\}$$

Claim C is r.e.

Claim $\bar{C}$ is not r.e. (hence C is not recursive)

Proof! Suppose C is r.e.
Let $\bar{C}$ be W_m

$$m \in C \iff \text{if } m \in W_m \iff m \in \bar{C}$$

Godel's Incompleteness Theorem

Lemma For every r.e. set A there is a formula $\varphi(x)$ of arithmetic such that $n \in A \iff \vdash \varphi(\bar{n})$

Hilbert 10th

$$x^2 + y^2 = 1$$

$$3xy^2 + 5x^3y - 13 = 0$$

Diophantine

$$P(x, y)$$

$$P(x_1, x_2, \dots, x_n) = 0$$

Doesn't have a solution in the integers

Let me go any $r \in \mathbb{N}$ set ~~the~~ we
can find a Diophantine eq $P(x, y_1, \dots, y_n)$

$n \in A \iff P(n, y_1, \dots, y_n) = 0$
has a solution in the integers

~~It is no~~ There is no algorithm
for deciding whether a Diophantine
equation ~~is~~ has a solution in
the integers or not

1/4 Col John Landry

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